## **Data Types**

Lecture II

#### Instructor: C. Pu (Ph.D., Assistant Professor)

puc@marshall.edu





- A *record* is an aggregate of data elements in which the individual elements are identified by *names* and accessed through offsets from the beginning of the structure.
- There is frequently a need in programs to model *a collection of data* in which the individual elements are *not of the same type or size*.
	- For example, information about a college student
		- Name, student number, grade point average
		- A character string for the name
		- An integer for the student number
		- A floating-point for the grade point average, and so forth
- *Records* are designed for this kind of need.





- It may appear that *records* and heterogeneous *arrays* are the same, but that is not the case.
- The elements of a heterogeneous *array* are all data objects with *same type*.
- The elements of a record are of potentially *different sizes* and reside in adjacent memory locations.





 In C, C++, and C#, *records* are supported with the *struct* data type.

```
struct structure name
{
  data_type member1;
  data_type member2;
   .
   .
  data_type memeber;
}
                                        struct student 
                                        {
                                             int idl;
                                             int id2;
                                             char a;
                                             char b;
                                             float percentage;
                                        };
```




 In C, C++, and C#, *records* are supported with the *struct* data type.





### **Record Types: Definitions of Records**

- The fundamental difference between a *record* and an *array* is that *record elements*, or *fields*, are *not referenced by index*.
- **Instead, the fields are named** with *identifiers*, and references to the fields are made using these *identifiers*.
- Another difference between *arrays* and *records* is that records in some languages are allowed to include *record* or *unions*.
- **The COBOL form of a record declaration, which is part of the data** division of a COBOL program, is illustrated in the following example:





### **Record Types: Definitions of Records**

- Ada uses a different syntax for records;
	- Rather than using the level numbers of COBOL, record structures are indicated by simply nesting record declarations inside record declarations.
- In Ada, records cannot be anonymous—they must be named types. Consider the followingAda declaration:

```
type Employee_Name_Type is record
            First : String (1..20);
            Middle : String (1..10);
            Last : String (1..20);
end record;
type Employee_Record_Type is record
            Employee_Name: Employee_Name_Type;
            Hourly_Rate: Float;
end record;
Employee_Record: Employee_Record_Type;
```


# **Record Types: Definitions of Records**

- In Java and C#, *records* can be defined as *data classes*, with nested records defined as nested classes.
- Data members of such classes serve as the record fields.

```
public class Student
{
  private String m_name;
  private int m_age;
  private String m_course;
  private String m_year;
  private String m_section;
  public Student( String name, int age, String course, String year, String section )
  {
     m name = name;m age = age;m course = course;m year = year;
     m section = section;
  }
  …
}
```




- References to the individual fields of records are syntactically specified by several different methods, two of which name the desired field and its enclosing records.
- COBOL field references have the form

*field\_name OF record\_name\_1 OF . . . OF record\_name\_n*

- The first record named is the *smallest or innermost record that contains the field*.
- The next record name in the sequence is that of the record that contains the previous record, and so forth.





 For example, the MIDDLE field in the COBOL record example above can be referenced with

#### **MIDDLE OF EMPLOYEE-NAME OF EMPLOYEE-RECORD**





- Most of the other languages use *dot notation* for field references
	- The components of the reference are connected with *periods*.
- Names in *dot notation* have the opposite order of COBOL references:
	- They use the name of the largest enclosing record first and the field name last.



 For example, the following is a reference to the field Middle in the earlier Ada record example:

#### **Employee\_Record.Employee\_Name.Middle**

```
type Employee_Name_Type is record
            First : String (1..20);
            Middle : String (1..10);
            Last : String (1..20);
end record;
type Employee_Record_Type is record
            Employee_Name: Employee_Name_Type;
            Hourly_Rate: Float;
end record;
Employee_Record: Employee_Record_Type;
```


 C and C++ use this same syntax for referencing the members of their structures.





## **Record Types: Implementation of Record Types**

- The fields of records are stored in *adjacent memory locations*.
- **But because the sizes of the fields are not necessarily the same,** the access method used for arrays is not used for records.
- **If a** Instead, the **offset address**, relative to the beginning of the record, is associated with each field.
- $\blacksquare$  Field accesses are all handled using these offsets.



- A *tuple* is a data type that is similar to a record, except that the elements are *not named*.
- Python includes an *immutable tuple* type.
	- If a tuple needs to be changed, it can be converted to an array with the list function.
	- After the change, it can be converted back to a tuple with the tuple function.
- One use of tuples is when an array must be *write protected*,
	- When it is sent as a parameter to an external function and the user does not want the function to be able to modify the parameter.



- Python's *tuples* are closely related to its *lists*, except that tuples are *immutable*.
- **A** tuple is created by assigning a tuple literal, as in the following example:

$$
myTuple = (3, 5.8, 'apple')
$$

- The elements of a tuple can be referenced with indexing in brackets, as in the following: **myTuple[1]**
	- This references the second element of the tuple, because tuple indexing begins at 0.



- Tuples can be contenated with the plus (*+*) operator.
- **They can be deleted with the del statement.**
- **There are also other operators and functions that operate on** tuples.



- ML includes a tuple data type.
- **An ML tuple must have at least two elements, whereas Python's** tuples can be empty or contain one element.
- **As in Python, an ML tuple can include elements of mixed types.**
- **The following statement creates a tuple:**

**val myTuple = (3, 5.8, 'apple');**





The syntax of a tuple element access is as follows:

#### **#1(myTuple);**

- This references the first element of the tuple.
- **A** new tuple type can be defined in ML with a type declaration, such as the following:

**type** intReal **= int \* real;**



## **List Types**

- *Lists* were first supported in the first functional programming language, LISP.
- **They have always been part of the functional languages, but in** recent years they have found their way into some imperative languages.
- **EXECU** Lists in Scheme and Common LISP are delimited by parentheses and the elements are not separated by any punctuation.
- **For example,**

## **(A B C D)**

**Nested lists have the same form, so we could have** 

#### **(A (B C) D)**



## **List Types**

- Data and code have the same syntactic form in LISP and its descendants.
	- If the list (A B C) is interpreted as code, it is a call to the function A with parameters B and C.
- **The fundamental list operations in Scheme are two functions** that take lists apart and that build lists.
- **The CAR** function returns the first element of its list parameter.
- $\blacksquare$  For example, consider the following example: **(CAR '(A B C))**





- The *CDR* function returns its parameter list minus its first element.
- **For example, consider the following example: (CDR '(A B C))**
	- $\blacksquare$  This function call returns the list (B C).



## **Union Types**

- A *union* is a type whose variables may store *different type* values at *different times* during program execution.
- Difference between *Struct* and *Union*





## **Union Types: Discriminated Versus Free Unions**

- C and C++ provide union constructs in which there is no language support for type checking.
- In C and C++, the *union* construct is used to specify union structures.
- **The unions in these languages are** called *free unions*, because programmers are allowed complete freedom from type checking in their use.
- **For example, consider the following C** union:

```
union flexType {
   int intEl;
   float floatEl;
};
union flexType el1;
float x;
. . .
ell.intEl = 27;
x = el l.floatEl;
```


### **Pointer and Reference Types**

- A *pointer* type is one in which the variables have a range of values that consists of *memory addresses*.
- Languages that provide a pointer type usually include two fundamental pointer operations: *assignment* and *dereferencing*.
- The first operation sets a pointer variable's value to some useful *address*.
- The second operation is used to access or manipulate data contained in memory location pointed to by a pointer.



## **Pointer and Reference Types: Pointer Operations**

- In C++, it is explicitly specified with the *asterisk* (*\**) as a prefix unary operator.
- **EXECONS** Consider the following example of dereferencing:
	- If ptr is a pointer variable with the value 7080 and the cell whose address is 7080 has the value 206, then the assignment





