Subprograms

Lecture 16

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- There are two distinct categories of subprograms—procedures and functions—both of which can be viewed as approaches to extending the language.
- All subprograms are collections of statements that define parameterized computations.
- **Functions** return values and **procedures** do not.
- In most languages that do not include procedures as a separate form of subprogram, functions can be defined not to return values and they can be used as procedures.



- Procedures can produce results in the calling program unit by two methods:
 - (1) If there are variables that are not formal parameters but are still visible in both the procedure and the calling program unit, the procedure can change them;
 - (2) If the procedure has formal parameters that allow the transfer of data to the caller, those parameters can be changed.



- Functions are called by appearances of their names in expressions, along with the required actual parameters.
- The value produced by a function's execution is returned to the calling code, effectively replacing the call itself.
- For example, the value of the expression f(x) is whatever value f produces when called with the parameter x.
- For a function that does not produce side effects, the returned value is its only effect.



- Functions define new user-defined "operators".
- For example,
 - if a language does not have an exponentiation operator, a function can be written that returns the value of one of its parameters raised to the power of another parameter.
- Its header in C++ could be

float power(float base, float exp)

which could be called with

result = 3.4 * power(10.0, x)



Local Referencing Environments: Local Variables

 In most contemporary languages, local variables in a subprogram are by default stack dynamic.

```
int adder(int list[], int listlen) {
    int sum = 0;
    int count;
    for (count = 0; count < listlen; count++)
        sum += list [count];
    return sum;
}</pre>
```



Local Referencing Environments: Local Variables

 In C and C++ functions, locals are stack dynamic unless specifically declared to be static.

```
int adder(int list[], int listlen) {
    static int sum = 0;
    int count;
    for (count = 0; count < listlen; count++)
        sum += list [count];
    return sum;
}</pre>
```



Local Referencing Environments: Local Variables

- Subprograms can define their own variables, thereby defining local referencing environments.
- Variables that are defined inside subprograms are called *local* variables, because their scope is usually the body of the subprogram in which they are defined.



Parameter-Passing Methods: Semantics Models of Parameter Passing

- Parameter-passing methods are the ways in which parameters are transmitted to and/or from called subprograms.
- Formal parameters are characterized by one of three distinct semantics models:
 - (I) They can receive data from the corresponding actual parameter;
 (*in mode*)
 - (2) They can transmit data to the actual parameter; (out mode)
 - (3) They can do both. (*inout mode*)



Parameter-Passing Methods: Semantics Models of Parameter Passing

- For example, consider a subprogram that takes two arrays of int values as parameters—list1 and list2.
 - The subprogram must add list1 to list2 and return the result as a revised version of list2.
 - Furthermore, the subprogram must create a new array from the two given arrays and return it.
 - For this subprogram, list1 should be in mode, because it is not to be changed by the subprogram.
 - list2 must be inout mode, because the subprogram needs the given value of the array and must return its new value.
 - The third array should be out mode, because there is no initial value for this array and its computed value must be returned to the caller.



Parameter-Passing Methods: Implementation Models of Parameter Passing

- A variety of models have been developed by language designers to guide the implementation of the three basic parameter transmission modes.
- The three semantics models of parameter passing when physical moves are used





Parameter-Passing Methods: Semantics Models of Parameter Passing

- There are two conceptual models of how data transfers take place in parameter transmission:
 - An *actual value* is copied (to the caller, to the called, or both ways),
 - An *access path* is transmitted.
- Most commonly, the access path is a simple pointer or reference.



Parameter-Passing Methods: Pass-by-Value

- When a parameter is *passed-by-value*, the value of the actual parameter is used to initialize the corresponding formal parameter, which then acts as a local variable in the subprogram, thus implementing in-mode semantics.
- Pass-by-value is normally implemented by copy, because accesses often are more efficient with this approach.





Parameter-Passing Methods: Pass-by-Value

Example:

```
void swap(int a, int b) {
  int temp;
 temp = a;
  a = b;
  b = temp;
int main() {
 int num I = 10, num 2 = 20;
 printf("Before swapping num I = \% d num 2 = \% d n", num I, num 2;
 swap(num1, num2);
 printf("After swapping num1 = %d num2 = %d\n", num1, num2);
 return 0;
```



Parameter-Passing Methods: Pass-by-Reference

- Pass-by-reference is a second implementation model for inoutmode parameters.
- Pass-by-reference method transmits an access path, usually just an address, to the called subprogram.
- This provides the access path to the cell storing the actual parameter.
 - Thus, the called subprogram is allowed to access the actual parameter in the calling program unit.
- In effect, the actual parameter is shared with the called subprogram.





Parameter-Passing Methods: Pass-by-Reference

Example:

```
void swap(int *a, int *b) {
  int temp;
  temp = *a;
  *a = *b:
  *b = temp;
int main() {
 int num I = 10, num 2 = 20;
 printf("Before swapping num1 = %d num2 = %d\n", num1, num2);
 swap(&num1, &num2);
 printf("After swapping num I = \%d num2 = \%d\n", num I, num2);
 return 0;
```

