# **Computer Networks and the Internet**

Instructor: C. Pu (Ph.D., Assistant Professor)

Lecture 02

puc@marshall.edu





## **Network Edge**

- Computers and other devices connected to the Internet = end systems
  - sit at the edge of the Internet
  - including
    - desktop computers
    - servers
    - mobile computers
    - more...
- end system also = hosts
  - because they host application programs
    - web browser program
    - web server program
    - email client and server program
- end system end system C host end system = host



## **Network Edge**

- Hosts can be further divided into
  - clients
    - desktop, mobile PC, smartphone, etc

client

CM

serve

- servers
  - more powerful machine that store and distribute information
- Most of the servers reside in large

### data centers





### **Network Core**

- Network core
  - the mesh of *packet switches* and *links* that interconnects the Internet's end systems
- In a network application
  - end systems exchange messages
  - messages can contain anything
    - perform control function
    - contain data







## **Network Core**

- To send a message from source to destination
  - source
    - breaks long messages into smaller chunks of data (*packets*)
  - between source and destination
    - each packet travels through communication links and packet switches







- Packet are transmitted over each communication link at a rate equal to the full transmission rate of the link
  - if a packet has L bits, transmitted over a link with transmission rate R bit/sec
  - then, the time to transmit the packet is L/R seconds







## **Store-and-Forward Transmission**

- Most packet switches use **Store-and-Forward Transmission** at the inputs to the links.
  - the packet switch must receive the entire packet before it can begin to transmit the first bit of the packet onto the outbound link



A router typically have many incident links, transferring a packet from one (incoming) link to one (outgoing) link.





- takes L/R seconds to transmit (push out) packet <a href="#"><u>Exar</u></a> of L bits on to link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- delay that destination receives the entire packet
  - 2L/R (assuming zero propagation delay)

### Example:

- L = 7.5 Mbits
- R = 1.5 Mbps
- transmission delay = 10 sec





- takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- store and forward: entire packet must arrive at router before it can be transmitted on next link
- delay to receive all three packets
  - 4L/R





 sending one packet from source to destination over a path consisting of N links each of rate R

• the end-to-end delay is: 
$$d_{\text{end-to-end}} = N \frac{L}{R}$$



# Packet Switching: Queueing Delay and Packet Loss

- Packet switch has multiple links
  - **output buffer** (also called an *output queue*) for each attached link
    - stores packets that the router is about to send into that link
  - If the link busy with the transmission of another packet?
    - packet must wait in the output buffer
- In addition to store-and-forward delay, there is an output buffer queuing delays
  - these delays are variable and depend on the level of congestion in the network
- The amount of buffer space is **limited** 
  - if the buffer is completely full when packet arrives
    - packet loss
    - either the arriving packet or one of the already-queued packets will be dropped





### queuing and loss:

- If arrival rate of packets exceeds transmission rate of link (1.5 Mb/s) for a period of time:
  - congestion will occur at the router as

**Packet Switching:** 

- packets will be queued in the buffer, wait to be transmitted on link
- packets can be dropped (lost) if memory (buffer) fills up



- A router takes a packet arriving on one of its attached communication links and forwards that packet onto another one of its attached communication links.
- But how does the router determine which link it should forward the packet onto?
  - Packet forwarding



- Every end system has an address called an IP address.
- When a source end system sends a packet to a destination end system
  - put destination's IP address in the packet's header
  - IP address "=" postal address
    - hierarchical structure
- When a packet arrives at a router
  - examines packet's destination address
  - forwards the packet to an adjacent router
    - forwarding table: maps destination addresses to that router's outbound links

When a packet arrives at a router, the router examines the address and searches its forwarding table, using this destination address, to find the appropriate outbound link.



- The end-to-end routing process is analogous to a car driver who does not use maps but instead prefers to ask for directions.
  - Suppose Joe is driving from New York City to 156 Lakeside Drive in Orlando (FL)
    - Joe first drives to point A, where the people tell Joe that he needs to get to point B.
    - So Joe drives to point B, where the people tell Joe that he needs to get to point C.
    - So Joe .....
    - So Joe drives to point x, where the people tell Joe that he just drives I miles east and he will reach 156 Lakeside Drive in Orlando (FL).



- How do forwarding tables get set?
- **routing protocols**: automatically set the forwarding tables
  - E.g., shortest path routing protocol
    - determine the shortest path from each router to each destination and use the shortest path results to configure the forwarding tables in the routers.



