Linked Lists

Lecture 03

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Adapted partially from Data Structures and Algorithms in Java, M.T. Goodrich, R.Tamassia and M. H. Goldwasser, Sixth Edition, Wiley; Data Structures and Algorithms in C++, Adam Drozdek, 4th Edition, Cengage Learning



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- Limitations of arrays,
 - the size of the array must be known at the time the code is compiled
 - the elements of the array are required potentially extensive shifting when inserting a new element
- linked lists, collections of
 - nodes storing data and links to other nodes
 - independent memory locations (nodes) that store data
 - links to other nodes
 - the addresses of the nodes
 - follow the links to move between nodes
- Utilize **pointer** to implement linked lists,
 - providing great flexibility

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- A node consists of **two** data members,
 - info store the node's information content
 - next point to the next node in the list



note: the next of last node is a null pointer.

- Each node is composed of ...
 - a data and a link (the address) to the next node in the sequence
- Use the single variable p (e.g., IntSLLNode *p) to access the entire list;
- Has a null pointer (\) in the last node in the list
- links have "direction"
- Let's create the linked list







IntSLLNode *p = new IntSLLNode(10);

- create the first node and make p point to this first node
- four steps:
 - a new IntSLLNode and p are created (Figure a)
 - set info of this new node to 10 (Figure b)
 - set next of this new node to null (Figure c)
 - make p point to this new node (Figure d)





- create the second node and make first node point to second node
- four steps:
 - a new IntSLLNode is created (Figure e)
 - set info of this new node to 8 (Figure f)
 - set next of this new node to null (Figure g)
 - make the first node point to this second node (Figure h)





- A disadvantage of single-linked lists:
 - the longer the list, the longer the chain of next pointers that need to be followed to a given node
 - reduce flexibility, and is prone to errors
- An alternative, use an additional pointer to the end of the list
 - keep two pointers: one to the first node; one to the last node



Singly Linked Lists (cont.) not part of the list; just for accessing the list (a) head tail (b) head tail (b) head tail (c) head tail (

- Uses two classes:
 - IntSLLNode, define the nodes of the list
 - IntSLList, define two pointers,
 - head and tail (e.g., IntSLLNode *head, *tail)





head = new IntSLLNode(el,head);

Insertion: a node is added at the <u>beginning</u> of a list

- . create a new (empty) node (figure a)
- 2. initialize the info member of the node (figure b)
- 3. initialize the next member to point to the first node in the list, which is the current value of *head* (figure c)
- 4. update the *head* to point to the new node (figure d)

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Insertion: a node at the end of a list

- create the new node and initialize the info member of the node (figures a and b)
- initialize the next member to null, since the node is at the end of the list (figure c)
- set the next member of the current last node to point to the new node (figure d)
- Since the new node is now the end of the list, update the tail pointer to point to it (figure e)
- if the list is initially empty, both head and tail would be set to point to the new node



```
else head = tail = new IntSLLNode(el);
```



- **Deletion**: at the <u>beginning</u> of the list
 - returning the value stored in the node
 - releasing the memory occupied by the node
 - first retrieve the value (info) stored in the first node
 - use a temporary pointer to point to the first node
 - set head to point to head \rightarrow next
 - delete the former first node, releasing its memory
 - note that,
 - when a single node is in the list, requiring that head and tail be set to null to indicate the list is now empty

- Deletion (cont.): at the end of a list
 - back the tail pointer to the previous node in list
 - this cannot be done directly
 - need a temporary pointer *tmp* to traverse the list until tmp → next = tail
 - once have located that node, retrieve the value contained in tail \rightarrow info, delete that node, and

```
set tail = tmp
int IntSLList::deleteFromTail() {
    int el = tail->info;
    if (head == tail) { // if only one node on the list;
        delete head;
        head = tail = 0;
    }
                         // if more than one node in the list,
    else {
        IntSLLNode *tmp; // find the predecessor of tail;
        for (tmp = head; tmp->next != tail; tmp = tmp->next);
        delete tail;
        tail = tmp;
                         // the predecessor of tail becomes tail;
        tail->next = 0;
    return el:
```



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