

# Sorting

## Lecture 19

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*Adapted partially from Data Structures and Algorithms in Java, M.T. Goodrich, R. Tamassia and M. H. Goldwasser, Sixth Edition, Wiley; Data Structures and Algorithms in C++, Adam Drozdek, 4th Edition, Cengage Learning*



# Introduction

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## ■ **Sorting**

- improve the **efficiency of accessing (or handling) data**, e.g., increasing or decreasing order; or alphabetical order
  - e.g, find a name in the telephone directory
    - alphabetically ordered
- *criteria* to order data: number or alphabetic character (ASCII)
- common/critical *properties of sorting algorithms* (machine-independent)
  - *number of comparisons*
  - *number of data movements*

} *compared and moved*

- may be difficult to determine exactly → approximations
- may differ depending on the original state of the data set (e.g., best case, worst case, and average case)

# Elementary Sorting Algorithms:

## Insertion Sort

- **Insertion sort:**

```
insertionsort(data[], n)
```

```
  for i = 1 to n-1 // unsorted set
```

```
    move all elements data[j] greater than data[i] by one position;  
    place data[i] in its proper position;
```

- The array of values → divide into two sets

- sorted values vs. unsorted

- Initially, the element with **index 0**

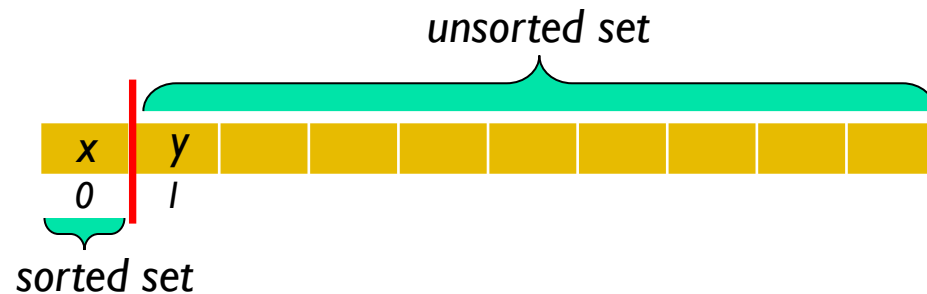
- sorted set

- The element with **index 1**

- the first element of the unsorted set

- Each repetition

- pick up the **first** element in the **unsorted set** and insert it into the correct **position of sorted set**



# Elementary Sorting Algorithms:

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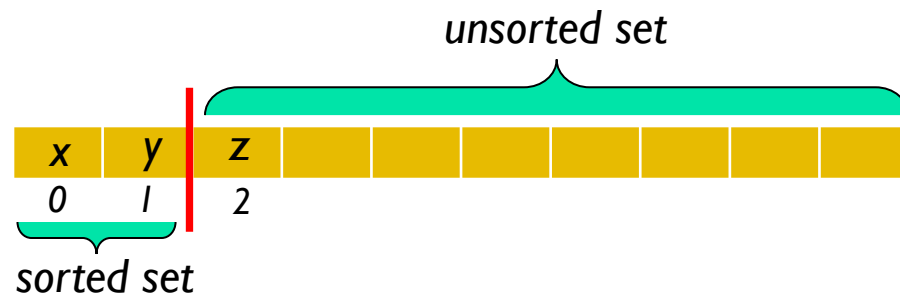
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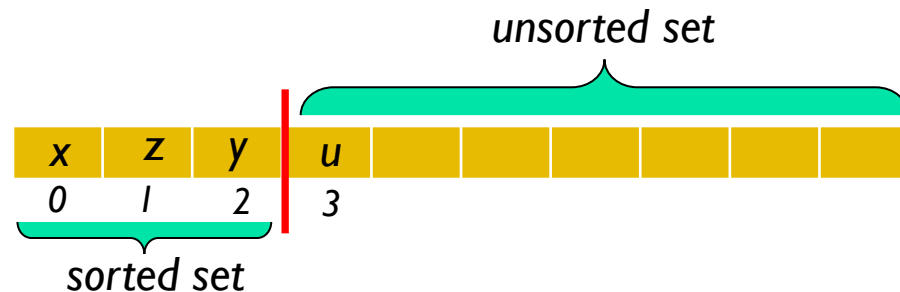
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# Elementary Sorting Algorithms:

## Insertion Sort (cont.)

- Consider an array of integers

39	9	45	63	18	81	108	54	72	36
----	---	----	----	----	----	-----	----	----	----

swap position



39	9	45	63	18	81	108	54	72	36
----	---	----	----	----	----	-----	----	----	----

$A[0]$  is the only element in sorted list

# Elementary Sorting Algorithms:

## Insertion Sort (cont.)

- The best case,

- the array is already **sorted**

1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10     *one compare*

1 | 2 | | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10     *one compare*

1 | 2 | 3 | | 4 | 5 | 6 | 7 | 8 | 9 | 10     *one compare*

- the first element in the unsorted set compares **only with the last element** of the sorted set
- $O(n)$

# Elementary Sorting Algorithms:

## Insertion Sort (cont.)

- The worst case,
  - the array is sorted in the **reverse order**

10 | 9 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1      *one compare*

9 | 10 | 8 | 7 | 6 | 5 | 4 | 3 | 2 | 1      *two compare*

8 | 9 | 10 | 7 | 6 | 5 | 4 | 3 | 2 | 1      *three compare*

- the first element in the unsorted set compares with **almost every element** in the sorted set
- $O(n^2)$





# Elementary Sorting Algorithms: Selection Sort

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- **Selection sort**
  - localize the exchange of array elements
    - finding a misplaced item and putting it in its final location
  - find the **first smallest value** in the array
    - place it in the **first** position
  - find the **second smallest value** in the array
    - place it in the **second** position
  - ...
  - repeat until the entire array is sorted

# Elementary Sorting Algorithms: Selection Sort (cont.)

- The pseudocode for the algorithm reflects its simplicity:

```
selectionsort (data [], n)
```

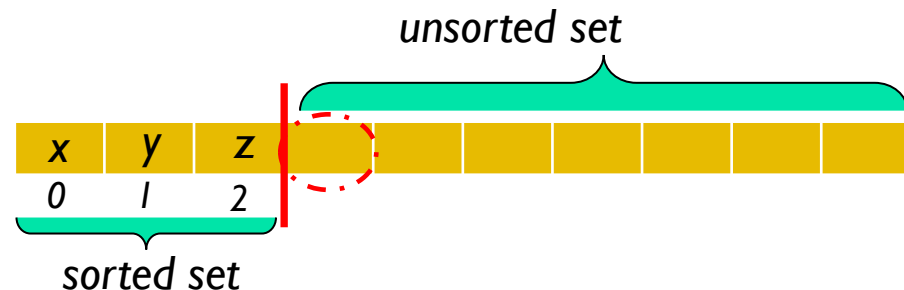
```
  for i = 0 to (n-2)
```

```
    select the smallest element among data[i], ..., data[n-1];
```

```
    swap it with data[i];
```

- Array divided into two sets:

- elements in the sorted set
- elements in the unsorted set



- The last value for  $i$  is  $n - 2$

- if all items have been looked at and placed except for the last item
  - the last element has to be the largest

# Elementary Sorting Algorithms: Selection Sort (cont.)

39	9	81	45	90	27	72	18
----	---	----	----	----	----	----	----

the location of smallest  
element in unsorted set



PASS	POS	ARR[0]	ARR[1]	ARR[2]	ARR[3]	ARR[4]	ARR[5]	ARR[6]	ARR[7]
------	-----	--------	--------	--------	--------	--------	--------	--------	--------



# Elementary Sorting Algorithms: Selection Sort (cont.)

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- Complexity
  - pass 1:
    - select the element with the first smallest value for all  $n$  elements
    - $n - 1$  comparisons
  - pass 2
    - select the element with the second smallest value for all  $n - 1$  elements
    - $n - 2$  comparisons
    - $(n - 1) + (n - 2) + \dots + 2 + 1$
    - $n(n - 1) / 2 = O(n^2)$



# Elementary Sorting Algorithms: Bubble Sort

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- **Bubble sort**
  - during each pass, **compare** pairs of adjacent items and **swap** them if they are in the wrong order
  - repeatedly moving the **largest (smallest) element** to the **highest (lowest) index position** of the array
  - continue till the list of unsorted elements exhaust
- The pseudocode of bubble sort:

```
bubblesort (data [], n)
  for i = 0 to n-2
    for j = n-1 down to i+1
      swap elements in positions j and j-1 if they are out of order;
```

# Elementary Sorting Algorithms:

## Bubble Sort (cont.)

- For example,

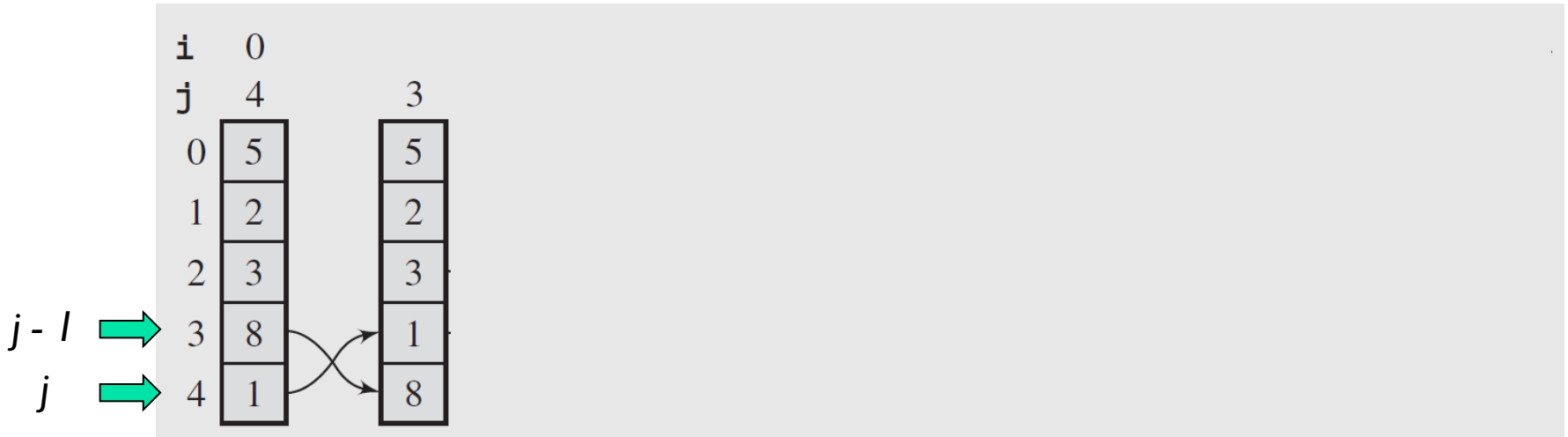
$i$	0
$j$	4
0	5
1	2
2	3
$j - 1$ →	3
$j$ →	4

Approach: repeatedly moving the smallest element to the lowest index position of the array.

# Elementary Sorting Algorithms:

## Bubble Sort (cont.)

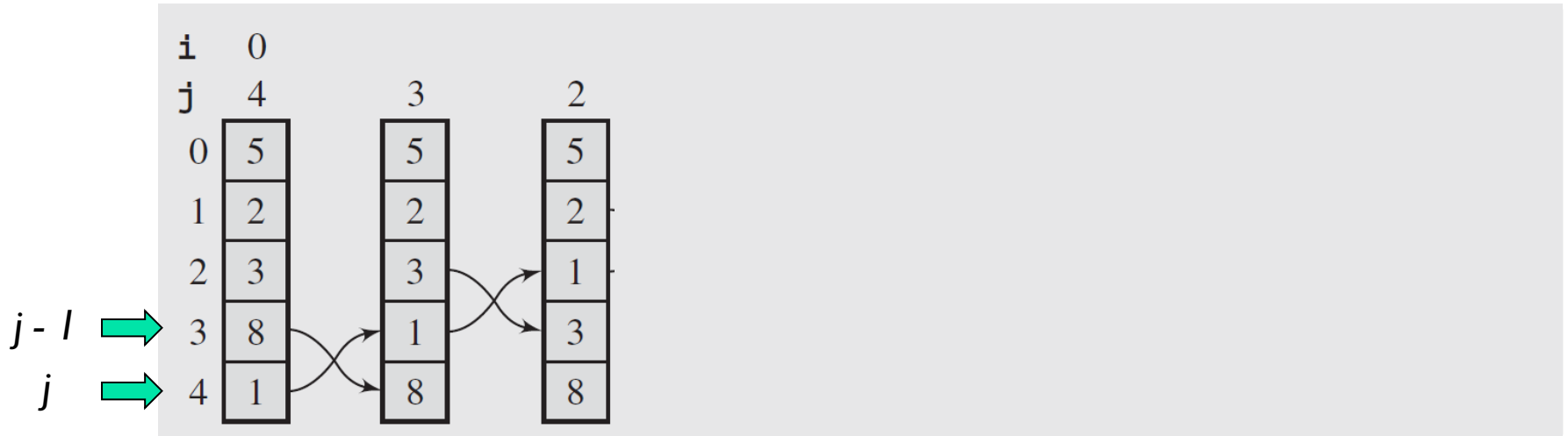
- For example,



# Elementary Sorting Algorithms:

## Bubble Sort (cont.)

- For example,

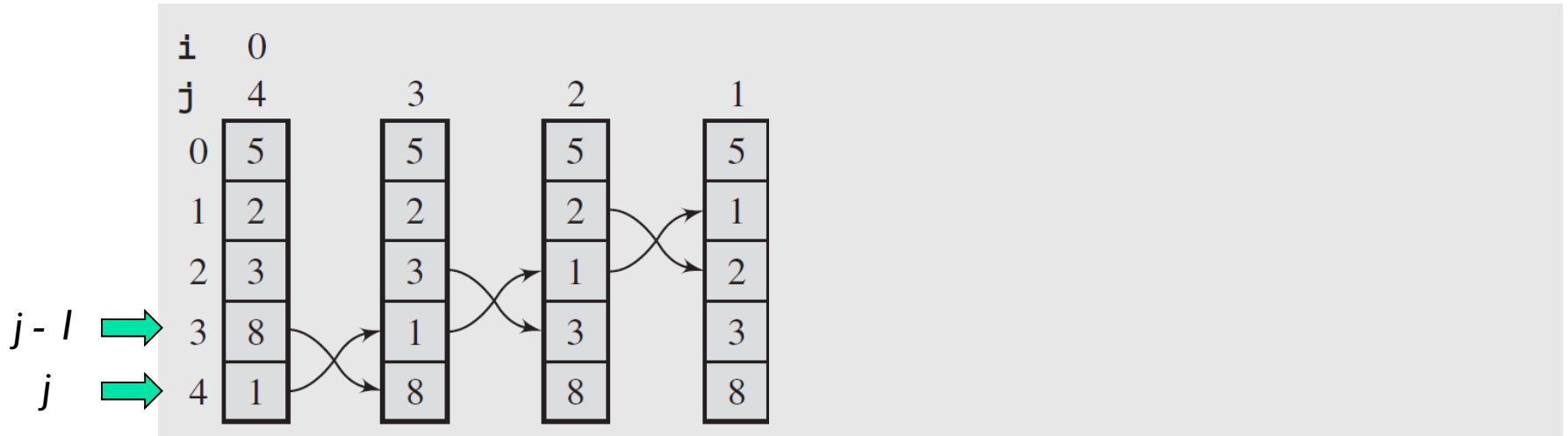




# Elementary Sorting Algorithms:

## Bubble Sort (cont.)

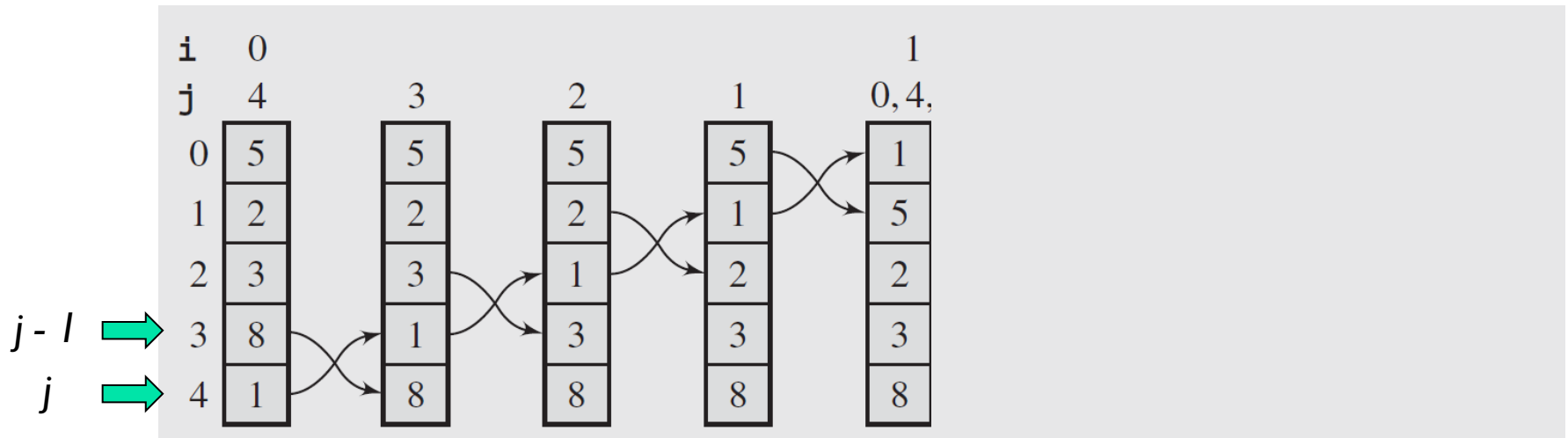
- For example,



# Elementary Sorting Algorithms:

## Bubble Sort (cont.)

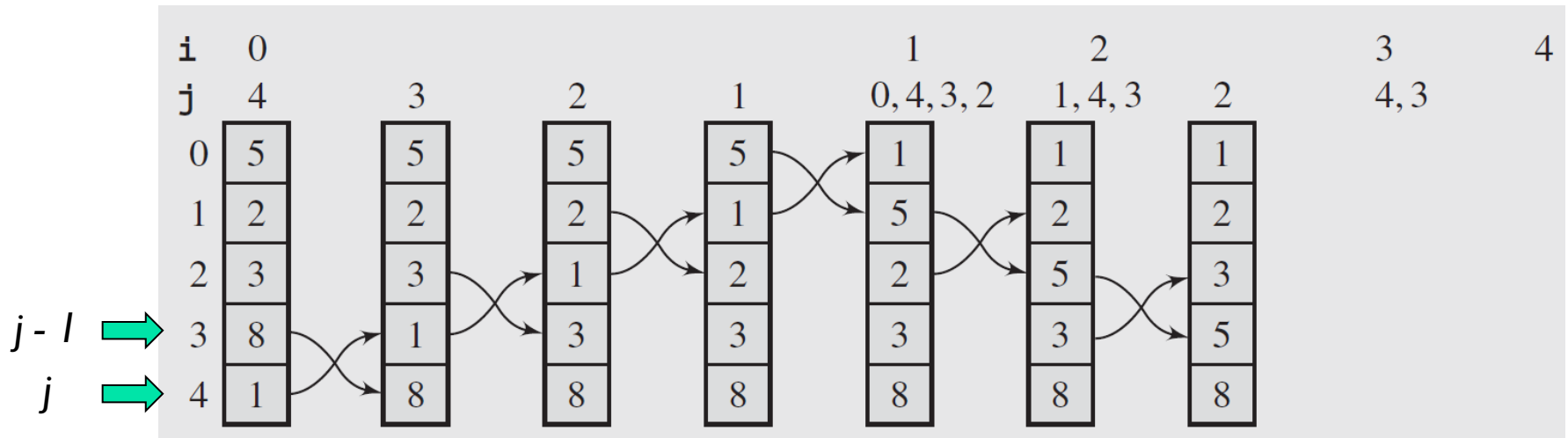
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# Elementary Sorting Algorithms:

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# Elementary Sorting Algorithms:

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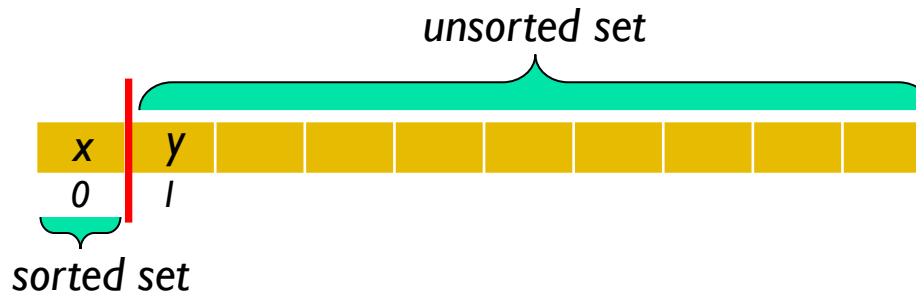
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- Complexity
  - in the first pass,
    - $n - 1$  comparisons
  - in the second pass,
    - $n - 2$  comparisons, and so on
  - $(n - 1) + (n - 2) + \dots + 2 + 1 = n(n - 1)/2 = O(n^2)$

# Efficient Sorting Algorithms:

## Shell Sort

- In insertion sort,
  - work well when the input element is “almost sorted”
  - if not sorted? inefficient to move the elements





# Efficient Sorting Algorithms: Shell Sort

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- **Shell sort**
  - improve over *insertion sort*
    - number of data movements
  - how?
    - sort parts (partial array) of the original array first and then;
    - if they are at least partially ordered or already sorted;
    - getting closer to the best case of an ordered array than initially.

# Efficient Sorting Algorithms:

## Shell Sort (cont.)

- The pseudo code of shell sort:

```
divide data into ( $h$  subarrays; ) one time or several times?  
for  $i = 1$  to  $h$   
    sort subarray  $data_i$ ;  
sort array data;
```

- if  $h$  is **too small**
  - the subarrays could be too large and the resulting sort would be inefficient
- if  $h$  is **too big**
  - too many subarrays
- use several different subdivisions
  - apply the same process separately to each subdivision

# Efficient Sorting Algorithms:

## Shell Sort (cont.)

- The pseudo code of shell sort (cont.):

```
determine numbers  $h_t \dots h_1$  of ways of dividing array data into subarrays;  
for (h= $h_t$ ; t > 1; t--, h= $h_t$ )  
    divide data into h subarrays;  
    for i = 1 to h  
        sort subarray  $data_i$ ;  
sort array data;
```

- called,
  - diminishing increment sort, shell sort, or shell's method



# Efficient Sorting Algorithms: Shell Sort (cont.)

- Perform the shell sort
  - arrange the elements in the form of a table
  - sort the columns
  - repeat with smaller number of long columns

63, 19, 7, 90, 81, 36, 54, 45, 72, 27, 22, 9, 41, 59, 33

$h = 8$

*Result:*

63	19	7	90	81	36	54	45	63	19	7	9	41	36	33	45
72	27	22	9	41	59	33		72	27	22	90	81	59	54	

63, 19, 7, 9, 41, 36, 33, 45, 72, 27, 22, 90, 81, 59, 54

$h = 5$

*Result:*

63	19	7	9	41	22	19	7	9	27
36	33	45	72	27	36	33	45	59	41
22	90	81	59	54	63	90	81	72	54

22, 19, 7, 9, 27, 36, 33, 45, 59, 41, 63, 90, 81, 72, 54

# Efficient Sorting Algorithms:

## Shell Sort (cont.)

- Perform the shell sort (cont.)
  - arrange the elements in the form of a table
  - sort the columns
  - repeat with smaller number of long columns

22, 19, 7, 9, 27, 36, 33, 45, 59, 41, 63, 90, 81, 72, 54

$h = 3$

22	19	7
9	27	36
33	45	59
41	63	90
81	72	54

*Result:*

9	19	7
22	27	36
33	45	54
41	63	59
81	72	90

9, 19, 7, 22, 27, 36, 33, 45, 54, 41, 63, 59, 81, 72, 90

# Efficient Sorting Algorithms:

## Shell Sort (cont.)

$h = 1$

*Result:*

9	7
19	9
7	19
22	22
27	27
36	33
33	36
45	41
54	45
41	54
63	59
59	63
81	72
72	81
90	90

- Perform the shell sort (cont.)
  - arrange the elements in the form of a table
  - sort the columns
  - repeat with smaller number of long columns

7, 9, 19, 22, 27, 33, 36, 41, 45, 54, 59, 63, 72, 81, 90

